# Good Approximations for the Relative Neighbourhood Graph 

Diogo Vieira Andrade (PUC-Rio)
Luiz Henrique de Figueiredo (IMPA)

## Outline

- Computational morphology
- The relative neighbourhood graph
- Computing the relative neighbourhood graph
- The Urquhart graph
- Results
- Conclusion
- Open problems


## Computational morphology

Computational morphology = computational extraction of perceptually meaningful structure from dot patterns.


Toussaint (1980) introduced RNG as tool for computational morphology.

## The relative neighbourhood graph

$S=$ set of points in the plane.
The edges in $\mathrm{RNG}(S)$ are defined by $p, q \in S$ with empty lune.


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## Computing the relative neighbourhood graph

- Brute-force algorithm from definition takes time $O\left(n^{3}\right)$.
- Restriction to $\mathrm{DT}(S)$ gives extraction in time $O\left(n^{2}\right)$.
- Supowit (1983) extracts in time $O(n \log n)$.
- Jaromczyk \& Kowaluk (1987) extract in time $O(n \alpha(n, n))$.
- Jaromczyk, Kowaluk \& Yao (1991?) extract in time $O(n)$.
- Lingas (1994) extracts in time $O(n)$
$\diamond$ simple algorithm, never implemented.


## The Urquhart graph

- Idea by Urquhart (1980): test only Delaunay neighbours!
$\diamond$ remove longest edge from each Delaunay triangle
$\diamond$ common mistake!
$\diamond$ new graph: Urquhart graph $\quad \mathrm{RNG}(S) \subseteq \mathrm{UG}(S) \subseteq \mathrm{GG}(S)$
- Toussaint (1980) proposed UG $(S)$ as approximation to $\operatorname{RNG}(S)$
- Our theme: how good is this approximation?
$\diamond$ How close is $\mathrm{UG}(S)$ to $\operatorname{RNG}(S)$ ?
- compare number of edges.
$\diamond$ Is UG $(S)$ good for computational morphology?
- see pictures!


## $U G \neq R N G$



## Results: random points in a square



Results: random points in a square


RNG 1241 edges
UG 1263 edges

Results: random points in a square


RNG 1241 edges
UG $1263=1241+22$ edges

Results: random points on a spiral

RNG


Results: random points on a spiral


RNG 1291 edges


UG 1301 edges

Results: random points on a spiral


RNG 1291 edges


UG $1301=1291+10$ edges

Results: random point on line art: earth


RNG


GG


Results: random point on line art: earth


RNG 1089 edges
UG 1116 edges

Results: random point on line art: earth


RNG 1089 edges
UG $\quad 1116=1089+27$ edges

Results: random point on line art: man


Results: random point on line art: man


RNG 663 edges
UG 682 edges

Results: random point on line art: man


RNG 663 edges


UG $682=663+19$ edges

## Conclusion

- UG $(S)$ good approximation to RNG $(S)$ :
$\diamond$ only about $2 \%$ additional edges for random samples
- Easy to extract $\mathrm{UG}(S)$ from $\mathrm{DT}(S)$ in linear time.
- Good, free, robust, optimal implementations of DT $(S)$ at netlib:
$\diamond$ Triangle, by Jonathan Richard Shewchuk
$\diamond$ sweep2, by Steve Fortune


## Open problems

- Compare implementations
$\diamond$ Supowit (1983)
$\diamond$ Lingas (1994)
- Probabilistic results à la Devroye (1988):
$\diamond E_{\mathrm{GG}}(N) \sim 2 N$
$\diamond E_{\mathrm{RNG}}(N) \sim(1.27+o(1)) N$
$\diamond E_{\mathrm{UG}}(N) \sim ? ? ? N$


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